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# **Target Applications**

- Fluid flow in porous media
  - Highly heterogeneous media
  - Interested in global properties of the solution
- Coupled fluid-elastic systems
  - Multiple material regimes
  - Different models require very different treatment
- Lattice quantum chromodynamics
  - Highly heterogeneous operator
  - Randomized heterogeneity within Monte Carlo process

### **Target Discrete Models**

- Finite difference/element models of elliptic systems
- Matrices are
  - sparse
  - symmetric
  - positive definite

Solving Ax = b

• Sparsity with large bandwidth means that direct methods are not effective for these problems

### **Algorithmic Goals**

Efficient and robust solver for heterogeneous models

Efficient: Optimal scalability, both algorithmic and parallel

- Cost of solve linearly proportional to number of unknowns
- Natural parallelism; most calculations should be data-local

Robust: Consistent performance with few parameters

- Predictable performance based on simple characteristics
- Not expert software; no magic parameters

# **Efficiency First**

Heterogeneity is an added complication, but not fundamental

Still need techniques to handle

- Large problem sizes
- Large condition numbers
- Multiscale structure of operator

Robustness without efficiency is EASY! Gaussian Elimination

# **Efficiency First**

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Robustness without efficiency is EASY! Gaussian Elimination

- Start with an efficient solver for homogeneous models
  - Geometric multigrid
- Look for where robustness and heterogeneity play a role

- Want to improve approximation,  $x^{(0)}$ , to  $x = A^{-1}b$
- Residual,  $r^{(0)}$ , is a measure of the error

$$r^{(0)} = b - Ax^{(0)} = Ax - Ax^{(0)} = A(x - x^{(0)})$$

• Choose 
$$M^{-1} pprox A^{-1}$$

• Take 
$$x^{(1)} = x^{(0)} + M^{-1}r^{(0)}$$

Error propagation form:  $e^{(1)} = (I - M^{-1}A)e^{(0)}$ 

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Error propagation form:  $e^{(1)} = (I - M^{-1}A)e^{(0)}$  $e^{(2)} = (I - M^{-1}A)^2 e^{(0)}$  $\vdots$  $e^{(n)} = (I - M^{-1}A)^n e^{(0)}$ 

### **Convergence of Stationary Iterations**



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- Instead of throwing out the method, look to complement its failings

How can we best correct error modes that are slow to be reduced by relaxation?

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- Instead of throwing out the method, look to complement its failings

How can we best correct error modes that are slow to be reduced by relaxation?

- Slow-to-converge errors are smooth
- Smooth vectors can be easily represented using fewer degrees of freedom

• Sine series representation:

$$f(x) = \sum_{k=1}^{\infty} c_k \sin(k\pi x)$$

• Discrete problems can only approximate certain modes

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Coarse grids accurately represent low-frequency modes Natural complement to relaxation

## Complementarity

Use two complementary processes to efficiently damp all errors

**Relaxation:** Damp high-frequency error by stationary iteration

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Use two complementary processes to efficiently damp all errors

**Relaxation:** Damp high-frequency error by stationary iteration

**Coarse-grid correction:** Eliminate low-frequency error by relaxation on coarse grids

Key realization: Solve for coarse-grid representation of error

• At any stage, error is reflected in residual:

$$r^{(k)} = b - Ax^{(k)} = Ax - Ax^{(k)} = A(x - x^{(k)})$$

• Don't transfer Ax = b to coarse grid, transfer Ae = r

## **The Details**

- Correct the approximation after relaxation,  $x^{(1)}$ , from an auxiliary (coarse-grid) problem
- Need interpolation map, P, from coarse grid to fine grid
- Corrected approximation will be  $x^{(2)} = x^{(1)} + Px_c$
- Optimal  $x_c$  satisfies  $P^T A P x_c = P^T r$ 
  - Optimal means  $||x x^{(2)}||_A$  is minimal

#### Multigrid choices:

- Coarse grid
- Interpolation operator, P
- Restriction, coarse-grid operator given by optimality

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#### Multigrid choices:

- Coarse grid
- Interpolation operator, P
- Restriction, coarse-grid operator given by optimality
- Restriction, coarse-grid operator given by physics
- Restriction, coarse-grid operator constrained by computation
#### **Multigrid Components** Relax: $x^{(1)} = x^{(0)} + D^{-1}r^{(0)}$

• Relaxation

- Use a smoothing process (such as Jacobi or Gauss-Seidel) to eliminate oscillatory errors
- Remaining error satisfies  $Ae^{(1)} = r^{(1)} = b Ax^{(1)}$



- Transfer residual to coarse grid
- Compute  $P^T r^{(1)}$



- Use coarse-grid correction to eliminate smooth errors
- Best correction,  $x_c$ , in terms of A-norm satisfies

$$P^T A P x_c = P^T r^{(1)}$$



Transfer correction to fine grid

• Compute 
$$x^{(2)} = x^{(1)} + Px_c$$



• Relax once again to remove oscillatory error introduced in coarse-grid correction



# Direct solution of coarse-grid problem isn't practical Recursion!

Apply same methodology to solve coarse-grid problem

# **Analysing Performance**

For simple problems, Fourier analysis predicts performance

Analysis based on invariant subspaces

- Eigenvectors of simple relaxation are Fourier modes
- Aliased pairs make invariant subspaces of interpolation and restriction
- Rediscretized CGO has same eigenvectors

Can directly compute convergence factor for two-level cycle

Fourier analysis only valid for operators with periodic character

Heterogeneity in A affects entire cycle

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- Convergence properties of relaxation change
- Slowest-to-converge errors reflect heterogeneity
- Interpolation must adapt to fit these errors
- Coarse-grid operator must account for heterogeneity

Must ensure complementarity is not lost

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Heterogeneity in the equations:

- coupled fluid-solid interactions, multiphysics
- blood flow in a vein, multi-species chemistry

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- blood flow in a vein, multi-species chemistry

Heterogeneity in the grid

- Local refinement, unstructured triangulation
- shock waves, irregular geometry

Aim for an algorithm that is robust to all of these

# **Multigrid Without Grids**

The essence of multigrid has nothing to do with grids!

Complementarity is key:

- Fix choice of relaxation
- For any A, some errors are slow to converge
- These errors must be corrected some other way

A. Brandt, S. McCormick, J. Ruge, in *Sparsity and Its Applications*, 1984 J. Ruge and K. Stüben, in *Multigrid Methods*, 1987

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Coarse-grid correction:

$$x \leftarrow x + PB_c^{-1}Rr$$
  
 $e \leftarrow e - PB_c^{-1}Rr$ 

A. Brandt, S. McCormick, J. Ruge, in *Sparsity and Its Applications*, 1984 J. Ruge and K. Stüben, in *Multigrid Methods*, 1987

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#### Variational Coarsening

Coarse-grid correction,

$$I - PB_c^{-1}RA,$$

can only correct errors in the range of P

Choosing  $R = P^T$  and  $B_c = P^T A P$  exactly eliminates errors in this space.

Complementarity is key:

- Errors reduced by relaxation and coarse-grid correction
- Errors that relaxation reduces slowly must be in range(P)

R. Nicolaides, Math. Comp. 1977, 31:892-906

# Algebraically Smooth Error

# Slow to converge errors of relaxation replace smooth modes within AMG

Design interpolation to accurately represent these modes

- Assume these errors give small residuals, Aepprox 0
- Expand residual equation:

$$a_{ii}e_i = -\sum_{j\in C}a_{ij}e_j - \sum_{k\notin C}a_{ik}e_k$$

• Use assumption on character of these errors to eliminate connections to  $k \notin C$ 

A. Brandt, S. McCormick, J. Ruge, in *Sparsity and Its Applications*, 1984 J. Ruge and K. Stüben, in *Multigrid Methods*, 1987

# **Graph-based coarsening**

Goal: Choose coarse-grid nodes to allow easy elimination of fine-fine connections

- Filter to eliminate small matrix entries
- Create graph of filtered matrix
- Greedy algorithm to choose maximal independent subset

Maximal independent subset ensures

- every fine-fine connection is "close" to a coarse-grid point
- coarse grid is small, but not too small

A. Brandt, S. McCormick, J. Ruge, in *Sparsity and Its Applications*, 1984 J. Ruge and K. Stüben, in *Multigrid Methods*, 1987

## **Big Picture Convergence**

Convergence is a spectral property

Write error-propagation operator as a matrix:

$$(I - M^{-1}A)(I - PB_c^{-1}P^TA)$$

Convergence factor is the spectral radius

$$\rho_{MG} = \lambda_{max} \left( (I - M^{-1}A)(I - PB_c^{-1}P^TA) \right)$$
  
or 
$$\rho_{MG} = \left\| (I - M^{-1}A)(I - PB_c^{-1}P^TA) \right\|_A$$

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**Expect:**  $\rho_{MG} = \sqrt{1 - \frac{1}{k}}$  for some  $k \ge 1$ **Goal:** Give an upper bound for k and, thus,  $\rho_{MG}$ 

J. Ruge and K. Stüben, in *Multigrid Methods*, 1987

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# Bounding $\rho_{MG}$

Abstract bounds on  $\rho_{\rm MG}$  come from interaction of relaxation and coarse-grid correction:

Define  $G = I - M^{-1}A$ ,  $T = I - P(P^T A P)^{-1}P^T A$ ,

Assume that  $\|Ge\|_A^2 \leq \|e\|_A^2 - \delta \|Te\|_A^2$ , then

 $\rho_{\rm MG} \leq \sqrt{1-\delta}$ 

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$$ho_{\mathsf{MG}} \leq \sqrt{1-\delta}$$

More commonly, we separate the assumptions:

Smoothing assumption:  $\|Ge\|_A^2 \leq \|e\|_A^2 - \alpha \|e\|_{AD^{-1}A}^2$ Approximation property:  $\|Te\|_A^2 \leq \beta \|Te\|_{AD^{-1}A}^2$ Then  $\rho_{MG} \leq \sqrt{1 - \frac{\alpha}{\beta}}$ 

J. Ruge and K. Stüben, in Multigrid Methods, 1987

#### **Smoothing Assumption**

Achieving bound on relaxation is "easy": Weighted Jacobi: For SPD A and  $\gamma_0 \ge \rho(D^{-1}A)$ , then if  $0 < \omega < \frac{2}{\gamma_0}$ ,  $\alpha \le \omega(2 - \omega\gamma_0)$ Gauss-Seidel: For SPD A,  $\alpha \le \frac{1}{(1+\gamma_0)(1+\gamma_0)}$ , for

 $\gamma_{-} = \sum_{j < i} rac{|a_{ij}|}{a_{ii}}$  and  $\gamma_{+} = \sum_{j > i} rac{|a_{ij}|}{a_{ii}}$ 

**Kaczmarz:** For any *A*, bound  $\alpha$  as in Gauss-Seidel using  $\gamma_{\pm}(AA^{T})$ 

A. Brandt, Appl. Math. Comput. 1986, 19:23-56

#### **Approximation Property**

Typically bound  $\beta$  indirectly:

$$\beta \leq \max_{e \neq 0} \min_{e_c} \frac{\|e - Pe_c\|_D^2}{\|e\|_A^2}$$

This leads to the Brandt-McCormick principle:

For each eigenvector, v, of A, interpolation must represent v with accuracy proportional to its eigenvalue

Challenge: How do we turn this into an algorithm?

- Partition grid
- Define interpolation coefficients

A. Brandt, Appl. Math. Comput. 1986, 19:23-56

S. McCormick and J. Ruge, SINUM 1985, 19:924-929

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#### **Generalized Approximation Property**

Instead of considering

$$\beta \leq \max_{e \neq 0} \min_{\mathbf{e}_c} \frac{\|\mathbf{e} - \mathbf{P} \mathbf{e}_c\|_D^2}{\|\mathbf{e}\|_A^2},$$

incorporate relaxation,  $I - M^{-1}A$ , into bound:

$$K = \max_{e \neq 0} \min_{e_c} \frac{\|e - Pe_c\|_{M(M+M^T-A)^{-1}M^T}^2}{\|e\|_A^2}$$

Then

$$ho_{\mathsf{MG}} \leq \left(1 - rac{1}{\mathcal{K}}
ight)^{rac{1}{2}}$$

R. Falgout and P. Vassilevski, SIAM J. Numer. Anal. 2004, **42**:1669-1693 Algebraic Multigrid Coarsening from theory to practice- p.25

# **Bounding** *K*

Let

• 
$$A = \begin{bmatrix} A_{ff} & -A_{fc} \\ -A_{cf} & A_{cc} \end{bmatrix}$$

• *M* be symmetric

• 
$$x^T A x \leq \omega x^T M x < 2 x^T M x$$

•  $\rho_f = \|I - M_{ff}^{-1} A_{ff}\|_{A_{ff}}$ 

Then,

$$\min_{P} \max_{e \neq 0} \min_{e_{c}} \frac{\|e - Pe_{c}\|_{M(M+M^{T}-A)^{-1}M^{T}}^{2}}{\|e\|_{A}^{2}} \leq \frac{1}{(2-\omega)(1-\rho_{f})}$$

For a given F/C partition, the best possible measure depends on the equivalence between  $M_{ff}$  and  $A_{ff}$ 

R. Falgout and P. Vassilevski, SIAM J. Numer. Anal. 2004, **42**:1669-1693 Algebraic Multigrid Coarsening from theory to practice- p.26

#### **Compatible Relaxation**

"A general measure for the quality of the set of coarse variables is the convergence rate of the compatible relaxation"

One approach:

- Run relaxation on tentative F-set
- Identify points where compatible relaxation is slow
- Choose subset of these points to add to C

This tells us about choosing C, but not about choosing P

A. Brandt, Elect. Trans. Numer. Anal. 2000, **10**:1-20
 O. Livne, Numer. Linear Algebra Appl. 2004, **2**:205-227

#### **Reduction-Based AMG**

Suppose we can partition the grid,  $\Omega = F \cup C$ , so that

$$x_f^T M_{ff} x_f \leq x_f^T A_{ff} x_f \leq \lambda_{\max} x_f^T M_{ff} x_f$$

and that  $\begin{bmatrix} M_{ff} & -A_{fc} \\ -A_{cf} & A_{cc} \end{bmatrix}$  is positive semi-definite. Choose **Relaxation:**  $I - \frac{2}{1+\lambda_{max}} \begin{bmatrix} M_{ff}^{-1} & 0 \\ 0 & 0 \end{bmatrix} A$ 

**Coarse-grid correction:** variational with  $P = \begin{bmatrix} M_{ff}^{-1}A_{fc} \\ I \end{bmatrix}$ 

Then

$$\rho_{\mathsf{MG}} \leq \left(1 - \left(\frac{2}{\lambda_{\mathsf{max}} + 1}\right)^2\right)^{\frac{1}{2}}$$

M. Ries, U. Trottenberg, G. Winter, J. Lin. Alg. Applic., 1983 S. MacLachlan, T. Manteuffel, S. McCormick, Numer. Linear Algebra Appl. 2006.

# **Putting It Together**

Goal: Use convergence theory to design algorithm **Pieces:** 

- Smoothing assumption
- Approximation property
- Generalized approximation property
- Compatible relaxation
  - Measure of quality of coarse-grid set
- AMGr
  - Choice of interpolation given good coarse-grid set

Still need an algorithm for partitioning

#### **Additive Multigrid**

Theory for additive preconditioners has same limitations.

Let  
• 
$$B = \begin{bmatrix} I & 0 \\ -A_{cf}M_{ff}^{-1}I \end{bmatrix} \begin{bmatrix} M_{ff} & 0 \\ 0 & S \end{bmatrix} \begin{bmatrix} I & -M_{ff}^{-1}A_{fc} \end{bmatrix}$$
  
•  $\begin{bmatrix} M_{ff} & -A_{fc} \\ -A_{cf} & A_{cc} \end{bmatrix}$  be positive semi-definite  
•  $x_{f}^{T}M_{ff}x_{f} \leq \lambda_{\min}x_{f}^{T}M_{ff}x_{f} \leq x_{f}^{T}A_{ff}x_{f} \leq \lambda_{\max}x_{f}^{T}M_{ff}x_{f}$   
•  $\nu_{\min}x_{c}^{T}Sx_{c} \leq x_{c}^{T}(A_{cc} - A_{cf}A_{ff}^{-1}A_{fc})x_{c} \leq \nu_{\max}x_{c}^{T}Sx_{c}$   
Then,

$$\kappa(B^{-rac{1}{2}}AB^{-rac{1}{2}}) \leq \left(1+\sqrt{1-rac{1}{\lambda_{\max}}}
ight)^2 rac{\lambda_{\max}^2 
u_{\max}}{\min(
u_{\min},\lambda_{\min})}.$$

O. Axelsson, Iterative Solution Methods, 1994

Y. Saad and B. Suchomel, Numer. Linear Algebra Appl. 2002, **9**:359-378

Y. Notay, Numer. Linear Algebra Appl. 2005, 12:419-451

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#### **Coarse-grid Selection**

Key to success in these bounds is in the partitioning of A

$$A = \left[ \begin{array}{cc} A_{ff} & -A_{fc} \\ -A_{cf} & A_{cc} \end{array} \right]$$

**Need**: Good approximation,  $M_{ff}$ , to  $A_{ff}$ **Need**: Cheap computation of  $M_{ff}^{-1}r_f$  and  $M_{ff}^{-1}A_{fc}$ **Need**: Dimension of  $A_{cc}$  much smaller than A

#### **Two Observations**

#### **1.** Cost of $M_{ff}^{-1}r_f$ depends on sparsity structure of $M_{ff}$

• Cheapest when  $M_{ff}$  is diagonal

S. MacLachlan, Y. Saad, SISC, to appear

#### **Two Observations**

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- **2.** Diagonally dominant  $A_{ff}$  can be approximated by its diagonal
  - More diagonally dominant  $\rightarrow$  better approximation
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  - $\blacktriangleright$  More diagonally dominant  $\rightarrow$  better approximation
- $A_{ff}$  is called  $\theta$ -dominant if, for each  $i \in F$ ,

$$a_{ii} \ge heta \sum_{j \in F} |a_{ij}|$$

S. MacLachlan, Y. Saad, SISC, to appear

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**Coarsening Goal:** Find largest set F such that  $A_{ff}$  is  $\theta$ -dominant.

S. MacLachlan, Y. Saad, SISC, to appear

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# **Optimality, Robustness, Complexity**

#### **Optimality:**

- Choosing diagonal  $M_{ff}$  ensures efficient relaxation
- Maximizing |F| so that  $A_{\rm ff}$  is  $\theta$ -dominant minimizes iteration cost

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#### **Robustness:**

• Fixing  $\theta$  guarantees  $\lambda_{\max} \leq \frac{1}{2\theta - 1}$ , bounds  $\rho_{MG}$ 

S. MacLachlan, Y. Saad, SISC, to appear

# **Optimality, Robustness, Complexity**

#### **Optimality:**

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**Robustness:** 

• Fixing  $\theta$  guarantees  $\lambda_{\max} \leq \frac{1}{2\theta-1}$ , bounds  $\rho_{MG}$ Complexity:

Finding max{|F| :  $A_{ff}$  is  $\theta$ -dominant}, is NP-complete.

- Constrained 0-1 integer programming problem
- Decision variables,  $f_i$ , are indicators of F/C partition
- Constraint function is row-wise  $\theta$ -dominance criterion

S. MacLachlan, Y. Saad, SISC, to appear

# **Greedy Algorithm**

Instead,

- Initialize  $U = \{1, \dots, n\}$ ,  $F = C = \emptyset$
- For each point in *U*, compute  $\hat{\theta}_i = \frac{a_{ii}}{\sum_{i \in F \cup U} |a_{ij}|}$

• Whenever 
$$\hat{\theta}_i \geq \theta$$
,  $i \to F$ 

- If  $U \neq \emptyset$ , then pick  $j = \operatorname{argmin}_{i \in U} \{ \hat{\theta}_i \}$ 
  - $j \rightarrow C$ • Update  $\hat{\theta}_i$  for all  $i \in U$  with  $a_{ii} \neq 0$

**Guarantee:**  $A_{ff}$  is  $\theta$ -dominant

Look for: largest  $A_{ff}$  possible in linear time

S. MacLachlan, Y. Saad, SISC, to appear

## A Challenge

AMGr theory requires  $P = \begin{bmatrix} M_{ff}^{-1}A_{fc} \\ I \end{bmatrix}$  and  $\begin{bmatrix} M_{ff} & -A_{fc} \\ -A_{cf} & Acc \end{bmatrix}$  be SPSD

${\cal K}$	grid	CA	$ ho_{MG}$
constant	$2048 \times 2048$	1.33	0.41
smooth	$2048 \times 2048$	1.33	0.41
binary	2048  imes 2048	1.98	0.91
anisotropic	32  imes 32	48.16	0.98989

Finite-element discretizations of  $-\nabla \mathcal{K} \nabla p$ 

In practice: Use greedy coarsening algorithm in combination with classical AMG interpolation

#### **Numerical Results**

$\mathcal{K}$	Grid	CA	ρ
constant	512  imes 512	1.33	0.13
	1024  imes 1024	1.33	0.14
	$2048 \times 2048$	1.33	0.14
smooth	512  imes 512	1.33	0.13
	1024  imes 1024	1.33	0.14
	$2048 \times 2048$	1.33	0.14
binary	512  imes 512	2.06	0.35
	1024  imes 1024	2.08	0.40
	$2048 \times 2048$	2.10	0.46
anisotropic	512  imes 512	2.39	0.13
	1024  imes 1024	2.41	0.20
	$2048 \times 2048$	2.43	0.20

#### **Time-Harmonic Maxwell Equations**

Maxwell Equations:

$$\begin{aligned} \nabla \times H &= J + \frac{\partial D}{\partial t}, & \nabla \cdot D &= \rho, \\ \nabla \times E &= -\frac{\partial B}{\partial t}, & \nabla \cdot B &= 0, \end{aligned}$$

Reduce system by assuming:

• linear constitutive laws

$$B = \mu H,$$
  $D = \epsilon E,$   $J = \sigma E$ 

- low-frequency excitation
- Constant 2D cross-section

#### **Induction Motor**

Find Fourier-domain potential,  $\hat{A} = (0, 0, \hat{A}_z)^T$ , by solving

$$-\nabla \cdot \left(\frac{1}{\mu} \nabla \hat{A}_z\right) + \imath \omega \sigma \hat{A}_z = \hat{J}_{s,z}$$



Lahaye et al., IEEE Trans. Magn. 2000, 36:1535-1538

#### **AMG** Performance



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# Summary

- Multigrid methods provide effective large-scale solvers
- Algebraic multigrid effective for heterogeneous systems
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  - Relaxation and coarse-grid correction

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- Multigrid methods provide effective large-scale solvers
- Algebraic multigrid effective for heterogeneous systems
- Complementarity is key
  - Relaxation and coarse-grid correction
- Convergence theory separates smoothing and approximation
- Compatible relaxation evaluates quality of coarse-grid set
- AMGr defines interpolation based on dominance principles
- Greedy algorithm puts these together

## **Limitations and Outlook**

Predictive AMG theory is very limited

AMGr & Greedy coarsening:

- In some cases, works as well as classical AMG
- Only guaranteed to be effective with diagonal dominance
   Goal: push both ways
- Find new algorithms based on old bounds
- Find new, predictive bounds on AMG theory