A parallel deflated Krylov solver for finite element problems

•

TU Delft

Kees Vuik, Guus Segal and Fred Vermolen c.vuik@math.tudelft.nl

http://ta.twi.tudelft.nl/users/vuik/

Delft University of Technology

Sparse Days and Grid Computing at St. Girons,

Hotel La Clairiere, St. Girons, France, June 10-13, 2003

C. Vuik, June 10, 2003 1 - p.1/?

Contents

•

- 1. Introduction
- 2. A parallel Krylov method for finite element problems
- 3. Deflation and Coarse Grid Acceleration
- 4. Numerical experiments
- 5. Conclusions

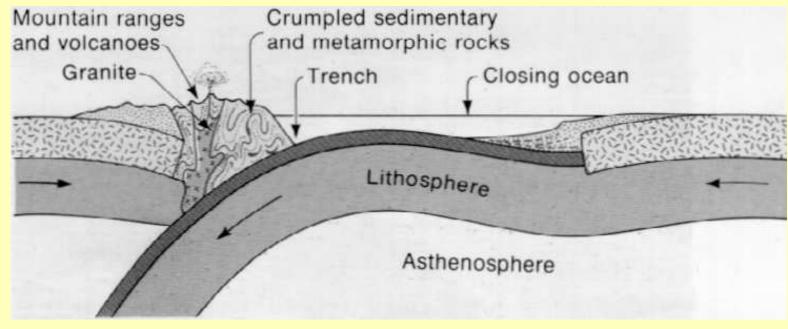
Reinhard Nabben, Jason Frank, Koos Meijerink, Erwin Dufour,

Gjalt Wijma, Larbi el Yaakoubi

1. Introduction

•

Motivation Knowledge of the fluid pressure in rock layers is important for an oil company to predict the presence of oil and gas in reservoirs.



The earth's crust has a layered structure

Incompressible Navier-Stokes problems

Discretized incompressible Navier-Stokes

- Momentum equations
- Pressure equation
- Transport equation

Coupled problem

$$\begin{pmatrix} \mathbf{Q} & \mathbf{G} \\ \mathbf{G}^T & \mathbf{0} \end{pmatrix} \begin{pmatrix} u \\ p \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \end{pmatrix}, \ u \in \mathbb{R}^n \text{ and } p \in \mathbb{R}^m$$

Solve the system Ax = b

TU Delft



•

Robust preconditioners

 (M)ICCG vd Vorst, Meijering, Gustafsson
 ILUT Saad, MRILU Ploeg, Wubs
 Navier-Stokes Elman, Silvester, Wathen, Golub
 RIF Benzi, Tuma



Literature review

- Robust preconditioners

 (M)ICCG vd Vorst, Meijering, Gustafsson
 ILUT Saad, MRILU Ploeg, Wubs
 Navier-Stokes Elman, Silvester, Wathen, Golub
 RIF Benzi, Tuma
- Parallel preconditioners Block variants see above ILU Bastian, Horton, Vuik, Nooyen, Wesseling SPAI Grote, Huckle, Benzi, Tuma, Chow, Saad

Literature review

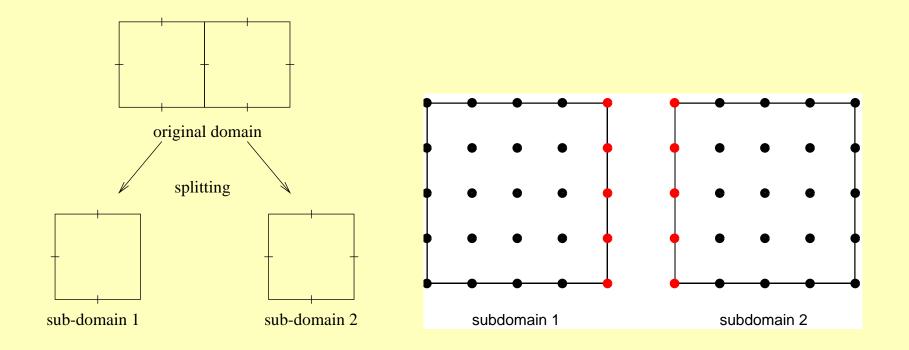
•

- Robust preconditioners

 (M)ICCG vd Vorst, Meijering, Gustafsson
 ILUT Saad, MRILU Ploeg, Wubs
 Navier-Stokes Elman, Silvester, Wathen, Golub
 RIF Benzi, Tuma
- Parallel preconditioners Block variants see above ILU Bastian, Horton, Vuik, Nooyen, Wesseling SPAI Grote, Huckle, Benzi, Tuma, Chow, Saad
- Acceleration of parallel preconditioners CGC Notay, vd Velde, Benzi, Frommer, Nabben, Szyld, Chan, Mathew, Dryja, Widlund, Padiy, Axelsson, Polman Deflation Nicolaides, Mansfield, Kolotilina, Frank, Vuik Morgan, Chapman, Saad, Burrage, Ehrel, Pohl FETI Farhat, Roux, Mandel, Klawonn, Widlund

2. A parallel Krylov method for finite element problems

Data distribution



TU Delft

Parallelization of ICCG

ICCG

$$\begin{split} k &= 0, \ r_0 = b - Ax_0, \ p_1 = z_1 = L^{-T}L^{-1}r_0;\\ \text{while } \|r_k\|_2 > \varepsilon \text{ do }\\ k &= k + 1;\\ \alpha_k &= \frac{(r_{k-1}, z_{k-1})}{(p_k, Ap_k)};\\ x_k &= x_{k-1} + \alpha_k p_k;\\ r_k &= r_{k-1} - \alpha_k Ap_k;\\ z_k &= L^{-T}L^{-1}r_k;\\ \beta_k &= \frac{(r_k, z_k)}{(r_{k-1}, z_{k-1})};\\ p_{k+1} &= z_k + \beta_k p_k; \end{split}$$

•

TU Delft

Explanation for a 1D example

Building blocks

•

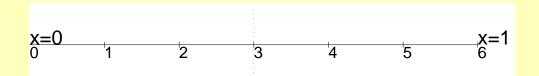
- vector update
- inner product
- matrix vector product
- preconditioner vector product

$$-\frac{d^2y}{dx^2} = f, \ y(0) = y(1) = 0.$$

Take n = 5 and decompose the domain into two subdomains (1 and 2)

Vector update

•



We define $I_1 = \{1, 2, 3, \}$ and $I_2 = \{3, 4, 5\}$. Note that there is an overlap of 1 point.

•

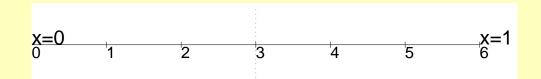
•

•

C. Vuik, June 10, 2003 9 - p.9/?

Vector update

•



We define $I_1 = \{1, 2, 3, \}$ and $I_2 = \{3, 4, 5\}$. Note that there is an overlap of 1 point.

Global vector
$$\mathbf{x} = \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \end{pmatrix}$$
, local vectors $\begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix}$ and $\begin{pmatrix} x_3 \\ x_4 \\ x_5 \end{pmatrix}$

Vector update is straight forward.

Inner product

•

- Determine the local innerproduct
- Sum the local innerproducts by MPI_ALLREDUCE



Inner product

- Determine the local innerproduct
- Sum the local innerproducts by MPI_ALLREDUCE

But



Inner product

- Determine the local innerproduct
- Sum the local innerproducts by MPI_ALLREDUCE

But

•

The contributions of the interface points are used more than once. Solution: use the interface points only in one local inner product.



C. Vuik, June 10, 2003 10 - p.10/?

Matrix vector product

•

•

TU Delft

Matrix vector product

•

$$A = \left(\begin{array}{cc} A_{11} & 0\\ 0 & 0 \end{array}\right) + \left(\begin{array}{cc} 0 & 0\\ 0 & A_{22} \end{array}\right)$$



Matrix vector product

•

$$A = \left(\begin{array}{cc} A_{11} & 0\\ 0 & 0 \end{array}\right) + \left(\begin{array}{cc} 0 & 0\\ 0 & A_{22} \end{array}\right)$$

The global matrix vector product $\mathbf{p} = A\mathbf{x}$:

1. Determine
$$\begin{pmatrix} p_1 \\ p_2 \\ p_3^l \end{pmatrix} = A_{11} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix}$$
 and $\begin{pmatrix} p_3^r \\ p_4 \\ p_5 \end{pmatrix} = A_{22} \begin{pmatrix} x_3 \\ x_4 \\ x_5 \end{pmatrix}$

in parallel.

- 2. Communication: send p_3^l from CPU1 to CPU2 and send p_3^r from CPU2 to CPU1. (nearest neighbour communication)
- 3. Determine on both processors $p_3 = p_3^l + p_3^r$ in parallel.

Parallelization of a block preconditioner

Take as preconditioner the following

$$\mathbf{p} = P^{-1}\mathbf{x} = \left(\sum_{i=1}^{p} R_i^T P_{i,i}^{-1} R_i\right) \mathbf{x}$$

where

•

$$P_{i,i} \approx A_{i,i}$$

C. Vuik, June 10, 2003 12 - p.12/?

Parallelization of a block preconditioner

Take as preconditioner the following

$$\mathbf{p} = P^{-1}\mathbf{x} = \left(\sum_{i=1}^{p} R_i^T P_{i,i}^{-1} R_i\right) \mathbf{x}$$

where

•

$$P_{i,i} \approx A_{i,i}$$

In our example

$$R_1 = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \end{pmatrix} \text{ and } R_2 = \begin{pmatrix} 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

Parallelization of a block preconditioner

The global preconditioner vector product $\mathbf{p} = P^{-1}\mathbf{x}$:

1. Determine
$$\begin{pmatrix} p_1 \\ p_2 \\ p_3^l \end{pmatrix} = P_{11}^{-1} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix}$$
 and $\begin{pmatrix} p_3^r \\ p_4 \\ p_5 \end{pmatrix} = P_{22}^{-1} \begin{pmatrix} x_3 \\ x_4 \\ x_5 \end{pmatrix}$ in parallel.

2. Communication: send p_3^l from CPU1 to CPU2 and send p_3^r from CPU2 to CPU1. (nearest neighbour communication)

C. Vuik, June 10, 2003 13 – p.13/?

3. Determine on both processors $p_3 = p_3^l + p_3^r$ in parallel.

TU Delft

3. Deflation and Coarse Grid Acceleration

A is SPD, Conjugate Gradients

 $P = I - AZE^{-1}Z^T$ with $E = Z^T AZ$

C. Vuik, June 10, 2003 14 - p.14/?

and $Z = [z_1...z_m]$, where $z_1, ..., z_m$ are independent deflation vectors.

Properties

•

- 1. $P^T Z = 0$ and P A Z = 0
- 2. $P^2 = P$
- 3. $AP^T = PA$

•

$$x = (I - P^T)x + P^T x,$$

TU Delft

• • • • • • • • •

•

$$x = (I - P^T)x + P^T x,$$

$$(I - P^T)x = ZE^{-1}Z^TAx = ZE^{-1}Z^Tb$$
,

•

$$x = (I - P^T)x + P^T x,$$

$$(I - \mathbf{P}^T)x = ZE^{-1}Z^TAx = ZE^{-1}Z^Tb, \qquad A\mathbf{P}^Tx = \mathbf{P}Ax = \mathbf{P}b.$$

•

$$x = (I - P^T)x + P^T x,$$

$$(I - P^T)x = ZE^{-1}Z^TAx = ZE^{-1}Z^Tb, \qquad AP^Tx = PAx = Pb.$$

$\begin{array}{ll} \textbf{DICCG} \\ k = 0, \ \hat{r}_0 = Pr_0, \ p_1 = z_1 = L^{-T}L^{-1}\hat{r}_0; \\ \textbf{while} \ \|\hat{r}_k\|_2 > \varepsilon \ \textbf{do} \\ k = k + 1; \\ \alpha_k = \frac{(\hat{r}_{k-1}, z_{k-1})}{(p_k, PAp_k)}; \\ x_k = x_{k-1} + \alpha_k p_k; \\ \hat{r}_k = \hat{r}_{k-1} - \alpha_k PAp_k; \\ z_k = L^{-T}L^{-1}\hat{r}_k; \\ \beta_k = \frac{(\hat{r}_k, z_k)}{(\hat{r}_{k-1}, z_{k-1})}; \end{array}$

end while TU Delft

Variants for values at interfaces

$$z_i=1$$
 on Ω_i and $z_i=0$ on $\Omega\setminus ar{\Omega}_i$

1. no overlap

•

- $z_i = 1$ at one subdomain
- $z_i = 0$ at other subdomains
- 2. complete overlap

 $z_i = 1$ at all subdomains

3. average overlap

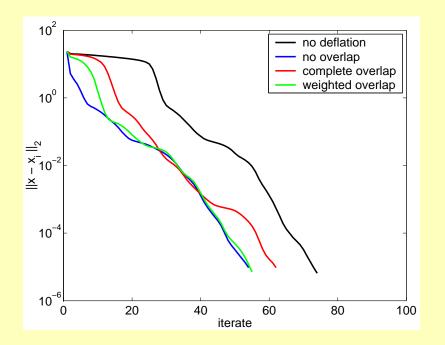
$$z_i = \frac{1}{n_{neighbors}}$$
 at all subdomains

4. weighted overlap ($-\operatorname{div}(\sigma \nabla u) = f$)

$$z_i = \frac{\sigma(i)}{\sum \sigma(neighbors)}$$

Error for Block IC and Deflation

Results for constant coefficients

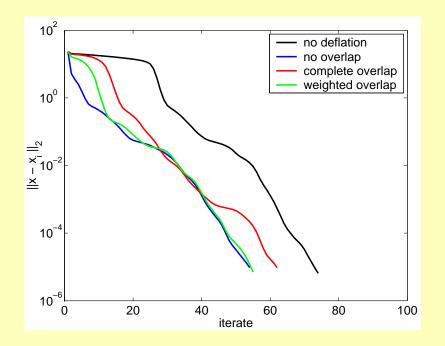


•

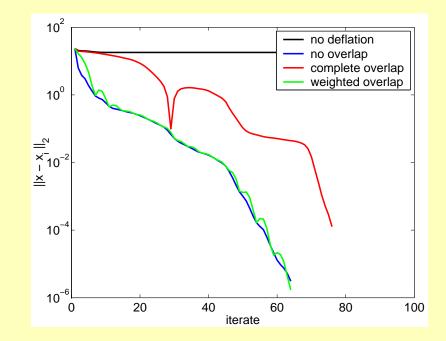
TU Delft

Error for Block IC and Deflation

Results for constant coefficients



and disontinuous coefficients



TU Delft

•

Parallel implementation (initialization)

Processor 1		Processor 2
Make z_1		Make z_2
	communication	
$z_{2\Gamma}$		$z_{1\Gamma}$
Make Az_1 and $Az_{2\Gamma}$		Make Az_2 and $Az_{1\Gamma}$
	communication	
sum up		sum up
$E_{11} = z_1^T A z_1$,		$E_{22} = z_2^T A z_2,$
$E_{12} = z_1^T A z_{2\Gamma}$		$E_{12} = z_2^T A z_{1\Gamma}$
	communication	
Determine Choleski		
decomposition of E		

TU Delft

•

Parallel implementation (during iteration)

$$P\mathbf{v} = \mathbf{v} - AZ(Z^T A Z)^{-1} Z^T \mathbf{v} = \mathbf{v} - AZE^{-1} Z^T \mathbf{v}$$

Processor 1		Processor 2
Compute $z_1^T v$		Compute $z_2^T v$
	communication	
$y = E^{-1} \left(\begin{array}{c} z_1^T v \\ z_2^T v \end{array} \right)$	communication	
$\mathbf{v} - y_1 A z_1 - y_2 A z_{2\Gamma}$		$\mathbf{v} - y_1 A z_{1\Gamma} - y_2 A z_2$

TU Delft

•

Coarse Grid Correction of ICCG

Definition

• •

- $Z \in \mathbb{R}^{n \times m}$ with independent columns.
- $E = Z^T A Z \in \mathbb{R}^{m \times m}$, E is SPD.
- $P_C = L^{-T}L^{-1} + \sigma Z E^{-1}Z^T$.

CICCG

W

$$\begin{split} k &= 0, \ r_0 = b - Ax_0, \ p_1 = z_1 = L^{-T}L^{-1}r_0; \\ \text{while} & \|r_k\|_2 > \varepsilon \text{ do} \\ & k = k + 1; \\ & \alpha_k = \frac{(r_{k-1}, z_{k-1})}{(p_k, Ap_k)}; \\ & x_k = x_{k-1} + \alpha_k p_k; \\ & r_k = r_{k-1} - \alpha_k Ap_k; \\ & z_k = P_C r_k = L^{-T}L^{-1}r_k + \sigma Z E^{-1}Z^T r_k; \\ & \beta_k = \frac{(r_k, z_k)}{(r_{k-1}, z_{k-1})}; \qquad p_{k+1} = z_k + \beta_k p_k \end{split}$$

end while

$$P_D = I - AZE^{-1}Z^T \qquad P_C = I + \sigma ZE^{-1}Z^T$$

TU Delft

•

 $P_D = I - AZE^{-1}Z^T \qquad P_C = I + \sigma ZE^{-1}Z^T$

Properties of P_D

- P_DA is symmetric and positive semidefinite
- P_D is a projection, $P_DAZ = 0$
- since P_DA is singular, a good termination criterion is important

 $P_D = I - AZE^{-1}Z^T \qquad P_C = I + \sigma ZE^{-1}Z^T$

Properties of P_D

- P_DA is symmetric and positive semidefinite
- P_D is a projection, $P_DAZ = 0$
- since P_DA is singular, a good termination criterion is important

C. Vuik, June 10, 2003 21 - p.21/?

Properties of P_C

- P_C is symmetric positive definite
- $A^{\frac{1}{2}}(P_C I)A^{\frac{1}{2}}$ is a projection

TU Delft

Definition Eigenpair $\{\lambda_i, v_i\}$, so $Av_i = \lambda_i v_i$ with $0 < \lambda_1 \leq \ldots \leq \lambda_n$. Take $Z = [v_1 \ldots v_m]$.

TU Delft

Properties of Deflation and CGC

Definition Eigenpair $\{\lambda_i, v_i\}$, so $Av_i = \lambda_i v_i$ with $0 < \lambda_1 \leq \ldots \leq \lambda_n$. Take $Z = [v_1 \ldots v_m]$.

Theorem

•

- the spectrum of $P_D A$ is $\{0, \ldots, 0, \lambda_{m+1}, \ldots, \lambda_n\}$
- the spectrum of $P_C A$ is $\{\sigma + \lambda_1, \ldots, \sigma + \lambda_m, \lambda_{m+1}, \ldots, \lambda_n\}$

Properties of Deflation and CGC

Definition Eigenpair $\{\lambda_i, v_i\}$, so $Av_i = \lambda_i v_i$ with $0 < \lambda_1 \leq \ldots \leq \lambda_n$. Take $Z = [v_1 \ldots v_m]$.

Theorem

•

- the spectrum of $P_D A$ is $\{0, \ldots, 0, \lambda_{m+1}, \ldots, \lambda_n\}$
- the spectrum of $P_C A$ is $\{\sigma + \lambda_1, \ldots, \sigma + \lambda_m, \lambda_{m+1}, \ldots, \lambda_n\}$

Corollary

DICCG converges faster than CICCG if $Z = [v_1 \dots v_m]$.

Notation: A, B are Hermitian, $A \succeq B$, if A - B is positive semidefinite



Notation: A, B are Hermitian, $A \succeq B$, if A - B is positive semidefinite

Some results from Horn and Johnson, Matrix Analysis

 $\lambda_k(A) + \lambda_1(B) \le \lambda_k(A+B) \le \lambda_k(A) + \lambda_n(B)$



Notation: A, B are Hermitian, $A \succeq B$, if A - B is positive semidefinite

Some results from Horn and Johnson, Matrix Analysis

 $\lambda_k(A) + \lambda_1(B) \le \lambda_k(A + B) \le \lambda_k(A) + \lambda_n(B)$

C. Vuik, June 10, 2003 23 - p.23/?

If A, B are positive definite with $A \succeq B$, then $\lambda_i(A) \ge \lambda_i(B)$.

TU Delft

Notation: A, B are Hermitian, $A \succeq B$, if A - B is positive semidefinite

Some results from Horn and Johnson, Matrix Analysis

 $\lambda_k(A) + \lambda_1(B) \le \lambda_k(A + B) \le \lambda_k(A) + \lambda_n(B)$

If A, B are positive definite with $A \succeq B$, then $\lambda_i(A) \ge \lambda_i(B)$.

Suppose that B has rank at most m. Then

- $\lambda_k(A+B) \leq \lambda_{k+m}(A), \quad k=1,2,\cdots,n-m,$
- $\lambda_k(A) \le \lambda_{k+m}(A+B), \quad k=1,2,\cdots,n-m.$

TU Delft

Deflation and Coarse Grid Correction (main result)

Theorem

•

Let A be symmetric positive definite and Z has rank Z = m. Let $E := Z^T A Z$. Then

C. Vuik, June 10, 2003 24 - p.24/?

$$\lambda_1(P_D A) = \dots = \lambda_m(P_D A) = 0$$

$$\lambda_n(P_D A) \leq \lambda_n(P_C A)$$

$$\lambda_{m+1}(P_D A) \geq \lambda_1(P_C A)$$

Deflation and Coarse Grid Correction (main result)

Theorem

•

Let A be symmetric positive definite and Z has rank Z = m. Let $E := Z^T A Z$. Then

$$\lambda_1(P_D A) = \dots = \lambda_m(P_D A) = 0$$

$$\lambda_n(P_D A) \leq \lambda_n(P_C A)$$

$$\lambda_{m+1}(P_D A) \geq \lambda_1(P_C A)$$

Theorem

 $Z_1 \in \mathbb{R}^{n \times r}$, $Z_2 \in \mathbb{R}^{n \times s}$, $rankZ_1 = r$ and $rankZ_2 = s$. If $ImZ_1 \subseteq ImZ_2$, then

$$\lambda_n((I - AZ_1E_1^{-1}Z_1^T)A) \geq \lambda_n((I - AZ_2E_2^{-1}Z_2^T)A)$$

$$\lambda_{r+1}((I - AZ_1E_1^{-1}Z_1^T)A) \leq \lambda_{s+1}((I - AZ_2E_2^{-1}Z_2^T)A)$$

Deflation and Coarse Grid Correction combined with a preconditioner

Definition

•

$$P_{CM^{-1}} := M^{-1} + \sigma Z E^{-1} Z^T.$$

Theorem

Let A and M be symmetric positive definite. Let $Z \in \mathbb{R}^{n \times m}$ with rankZ = m. Let $E := Z^T A Z$. Then

$$\lambda_n(M^{-1}P_DA) \leq \lambda_n(P_{CM^{-1}}A),$$

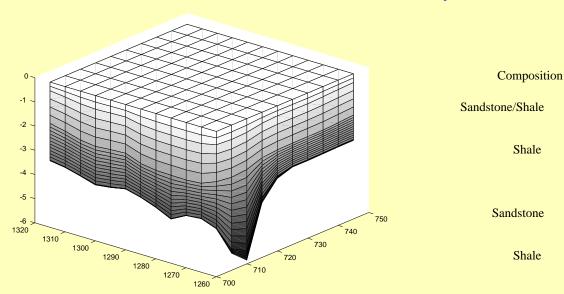
$$\lambda_{m+1}(M^{-1}P_DA) \geq \lambda_1(P_{CM^{-1}}A).$$

Corollary

DICCG converges faster than CICCG for general projection vectors.

C. Vuik, June 10, 2003 25 - p.25/?

4. Numerical experiments



Oil flow problem

method	Deflation	CGC
iterations	36	47
CPU time	5.9	8.2

TU Delft

•

Permeability

-4 10

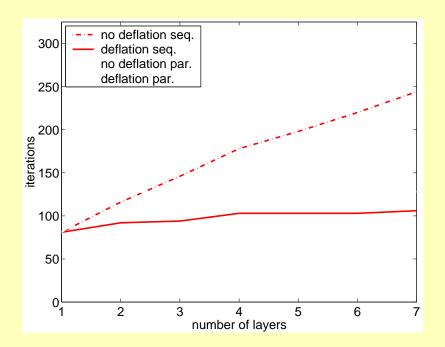
-7 10

10

-7 10

Poisson on parallel layers

Iterations

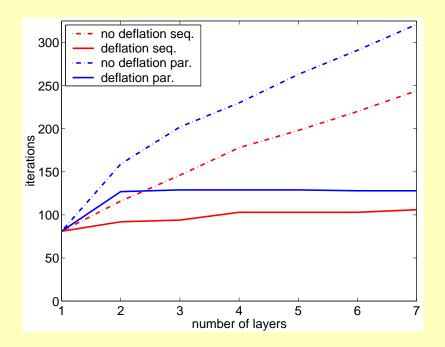


•

TU Delft

Poisson on parallel layers

Iterations

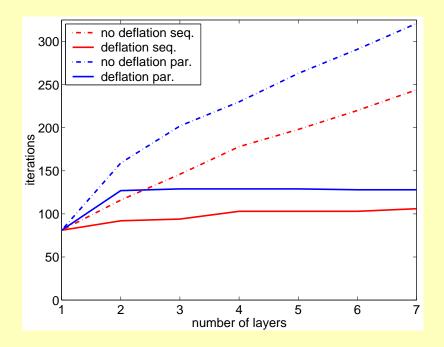


•

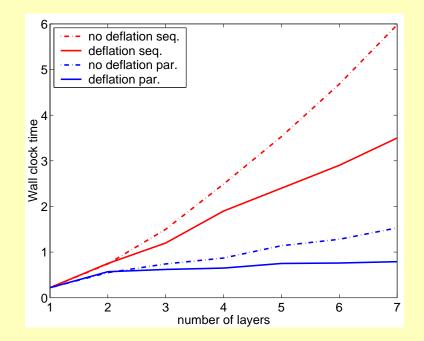
TU Delft

Poisson on parallel layers

Iterations



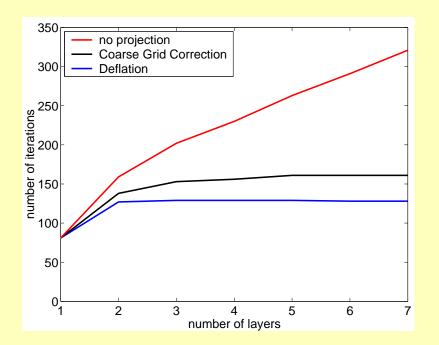
Wall clock time



TU Delft

Poisson on layers and blocks

Layers

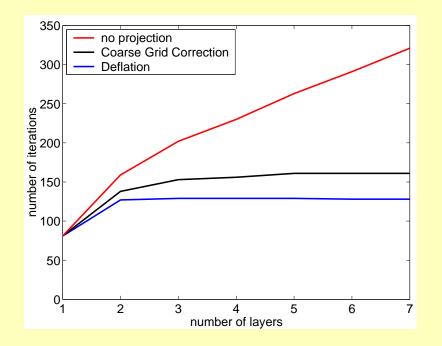


•

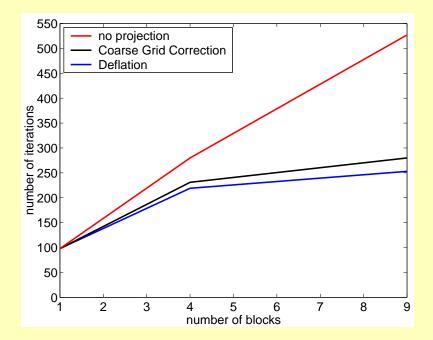
TU Delft

Poisson on layers and blocks

Layers



•



Blocks

TU Delft

5. Conclusions

- Block preconditioned Krylov methods combined with Deflation or CGC are well parallelizable (scalable, good speed up).
- For the vertex centered case, the weighted overlap strategy is optimal
- DICCG is more efficient than CICCG.
- Choices for the deflation vectors lead to comparable results in DICCG and CICCG.
- DICCG is a robust and efficient method to solve diffusion problems with discontinuous coefficients.



Further information

•

- http://ta.twi.tudelft.nl/nw/users/vuik/pub_it_def.html
- C. Vuik, A. Segal and J.A. Meijerink
 J. Comp. Phys., 152, pp. 385-403, 1999.
- J. Frank and C. Vuik SIAM Journal on Scientific Computing, 23, pp. 442–462, 2001
- C. Vuik, A. Segal, L. El Yaakoubi and E. Dufour Applied Numerical Mathematics, 41, pp. 219–233, 2002
- F.J. Vermolen, C. Vuik and A. Segal J. of Comp. Methods in Sciences and Engineering, to appear
- R. Nabben and C. Vuik A comparison of Deflation and Coarse Grid Correction, to appear

Overview

•

Krylov	Ar
Preconditioned Krylov	$L^{-T}L^{-1}Ar$
Block Preconditioned Krylov	$\sum_{i=1}^m (L_i^{-T}L_i^{-1})Ar$
Block Preconditioned Deflated Krylov	$\sum_{i=1}^{m} (L_i^{-T} L_i^{-1}) PAr$

TU Delft

•

•